The answers must be the original work of the author. While discussion with others is permitted and encouraged, the final work should be done individually. You are not allowed to work in groups. You are allowed to build on the material supplied in the class. Any other source must be specified clearly.

**1.** (60 points) Give a context-free grammar for each of the following languages. Explain how the grammar works.

No points will be given if the CFG is not commented.

- (a)  $L = \emptyset$
- **(b)**  $L = \{\lambda\}$
- (c)  $L = \{w \mid w \in \{a, b, c\}^*, w \text{ does not contain substring } bc\} = c^*(b \cup ac^*)^*$
- (d) ROLL is the language generated by the following recursive definition. It gives a few basic drum roll patterns. R stands for "right" and L stands for "left."

**basis:** "single-stroke:"  $\in$  ROLL, "double-stroke:"  $\in$  ROLL, "silly-stroke:"  $\in$  ROLL.

**recursive step:** If  $(w \in ROLL \text{ and } w \text{ contains "single"}) \text{ then } wRL \text{ is in } ROLL$ 

If  $(w \in ROLL \text{ and } w \text{ contains "double"}) \text{ then } wRRLL \text{ is in } ROLL$ 

If  $(w \in ROLL \text{ and } w = xy \text{ where } x = \text{"silly-stroke:"}) \text{ then } xRyL \text{ is in } ROLL$ 

**closure**: A string  $w \in \text{ROLL}$  only if it can be obtained from the basis set by a finite number of applications of the recursive step.

- (e)  $L = \{a^n b^m c^{2n+m} \mid n, m \ge 0\}$
- (f)  $L = \{a^n b^m \mid n \neq m, n \geq 0, m \geq 0\}$

(Hint: "not equal to" means "less than or greater than".)

**2.** (20 + 20 points) Consider the following grammar:

$$S \to aSA \mid \lambda$$
$$A \to bA \mid \lambda$$

- (a) Prove the grammar is ambiguous by finding a string that has two distinct leftmost derivations. Show the two derivations.
- (b) Build the derivation trees for the derivations in Part (a).